

STATS 217: Introduction to Stochastic Processes I

Lecture 18

Coupling and total variation

(μ, ν) $(\tilde{\mu}, \tilde{\nu})$

$$TV(\mu, \nu) = \frac{1}{2} \sum_x |\mu(x) - \nu(x)|$$

- Let μ and ν be two probability distributions on Ω .
- The **coupling lemma** asserts that (very fundamental/useful).

$$TV(\mu, \nu) = \inf\{\mathbb{P}[X \neq Y] : (X, Y) \text{ is a coupling of } \mu \text{ and } \nu\}.$$

extreme case: $\mu = \nu$

$$TV(\mu, \mu) = 0$$

a coupling of μ & μ is
just (X, X)
where $X \sim \mu$.

last time:

$$\mu = \text{Ber}(p)$$

$$\nu = \text{Ber}(q)$$

↳ ind. coupling $\frac{\mathbb{P}[X \neq Y]}{2}$
 ↳ monotone coupling $\mathbb{P}[X \neq Y]$
 ↳ ...

* if you give me any coupling \rightarrow upper bound on TV.
 (this is what we will use for e.g. Poisson approx.)

* if you can compute TV, 'test' quality of a coupling.

Coupling and total variation

- Let μ and ν be two probability distributions on Ω .
- The **coupling lemma** asserts that

$$\text{TV}(\mu, \nu) = \inf\{\mathbb{P}[X \neq Y] : (X, Y) \text{ is a coupling of } \mu \text{ and } \nu\}.$$

- **Example:** let $\mu = \text{Ber}(p)$ and $\nu = \text{Ber}(q)$ with $0 \leq p \leq q \leq 1$.
- Then, by direct computation,

$$\text{TV}(\text{Ber}(p), \text{Ber}(q)) = \frac{1}{2} (|q - p| + |1 - q - 1 + p|) = \bar{q} - p.$$

$$\frac{1}{2} |\mu(0) - \nu(0)| + \frac{1}{2} |\mu(1) - \nu(1)|$$

Coupling and total variation

- Let μ and ν be two probability distributions on Ω .
- The **coupling lemma** asserts that

$$\text{TV}(\mu, \nu) = \inf\{\mathbb{P}[X \neq Y] : (X, Y) \text{ is a coupling of } \mu \text{ and } \nu\}.$$

(intuitively, same as min)

- **Example:** let $\mu = \text{Ber}(p)$ and $\nu = \text{Ber}(q)$ with $0 \leq p \leq q \leq 1$.
- Then, by direct computation,

$$\text{TV}(\text{Ber}(p), \text{Ber}(q)) = \frac{1}{2}(|q - p| + |1 - q - 1 + p|) = q - p.$$

- For the monotone coupling (\hat{X}, \hat{Y}) , we have

$$\mathbb{P}[\hat{X} \neq \hat{Y}] = \mathbb{P}[1 - q \leq r \leq 1 - p] = q - p.$$

$\hat{x} = \hat{y} = 0$ dis. $\hat{x} = \hat{y} = 1$

\Rightarrow monotone coupling is optimal.

Coupling and total variation

- Let μ and ν be two probability distributions on Ω .
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- The above characterization shows that the monotone coupling is an **optimal coupling**.

Coupling and total variation

- $TV(\mu, \nu) = \overline{\inf} \{ \mathbb{P}[X \neq Y] : (X, Y) \text{ is a coupling of } \mu \text{ and } \nu \}$.
- Easy/useful direction: \leq . Why?

it suffices to show that for any coupling
 (X, Y) of μ & ν , we have

$$TV(\mu, \nu) \leq \mathbb{P}[X \neq Y]$$

& now take inf on both sides
OVER all couplings

Coupling and total variation

- $TV(\mu, \nu) = \inf\{\mathbb{P}[X \neq Y] : (X, Y) \text{ is a coupling of } \mu \text{ and } \nu\}$.
- Easy/useful direction: \leq . Why?
- Let (X, Y) be any coupling of μ, ν . Let $A \subseteq \Omega$.

$$TV(\mu, \nu) = \sup_{A \subseteq \Omega} |\mu(A) - \nu(A)|$$

enough to show that given $A \subseteq \Omega$

$$|\mu(A) - \nu(A)| \leq \mathbb{P}(X \neq Y)$$

and now, take max over A on both sides

Coupling and total variation

- $\text{TV}(\mu, \nu) = \inf\{\mathbb{P}[X \neq Y] : (X, Y) \text{ is a coupling of } \mu \text{ and } \nu\}$.
- Easy/useful direction: \leq . Why?
- Let (X, Y) be any coupling of μ, ν . Let $A \subseteq \Omega$. Then,

$$\underline{\underline{\mu(A)}} - \underbrace{\nu(A)} = \underline{\underline{\mathbb{P}[X \in A]}} - \underbrace{\mathbb{P}[Y \in A]}$$

$$\begin{aligned} X &\sim \mu \\ Y &\sim \nu \end{aligned}$$

Coupling and total variation

- $TV(\mu, \nu) = \inf \{ \mathbb{P}[X \neq Y] : (X, Y) \text{ is a coupling of } \mu \text{ and } \nu \}$.
- Easy/useful direction: \leq . Why?
- Let (X, Y) be any coupling of μ, ν . Let $A \subseteq \Omega$. Then,

$$\begin{aligned} \mu(A) - \nu(A) &= \mathbb{P}[X \in A] - \mathbb{P}[Y \in A] \\ &= \mathbb{P}[X \in A] - \mathbb{P}[X \in A, Y \in A] - \mathbb{P}[X \notin A, Y \in A] \\ &\quad \underbrace{\hspace{1.5cm}}_{\text{wavy}} \quad \underbrace{\hspace{1.5cm}}_{\text{underline}} \quad \underbrace{\hspace{1.5cm}}_{\text{underline}} \\ &\quad \underbrace{\hspace{1.5cm}}_{\text{double underline}} \quad \underbrace{\hspace{1.5cm}}_{\text{double underline}} \\ &= \mathbb{P}[X \in A, Y \notin A] + \mathbb{P}[X \in A, Y \in A] - \mathbb{P}[X \in A, Y \in A] - \mathbb{P}[X \notin A, Y \in A] \\ &= \mathbb{P}[X \in A, Y \notin A] - \mathbb{P}[X \notin A, Y \in A] \end{aligned}$$

Coupling and total variation

- $TV(\mu, \nu) = \inf\{\mathbb{P}[X \neq Y] : (X, Y) \text{ is a coupling of } \mu \text{ and } \nu\}$.
- Easy/useful direction: \leq . Why?
- Let (X, Y) be any coupling of μ, ν . Let $A \subseteq \Omega$. Then,

$$\begin{aligned}\mu(A) - \nu(A) &= \mathbb{P}[X \in A] - \mathbb{P}[Y \in A] \\ &= \mathbb{P}[X \in A] - \mathbb{P}[X \in A, Y \in A] - \mathbb{P}[X \notin A, Y \in A] \\ &= \mathbb{P}[X \in A, Y \notin A] - \mathbb{P}[X \notin A, Y \in A]\end{aligned}$$

$$\leq \mathbb{P}[X \in A, Y \notin A]$$

$$\leq \mathbb{P}[X \neq Y]$$

$$\begin{aligned}\{X \in A, Y \notin A\} \\ \subseteq \{X \neq Y\}\end{aligned}$$

Coupling and total variation

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- Easy/useful direction: \leq . Why?
- Let (X, Y) be any coupling of μ, ν . Let $A \subseteq \Omega$. Then,

$$\begin{aligned}\mu(A) - \nu(A) &= \mathbb{P}[X \in A] - \mathbb{P}[Y \in A] \\ &= \mathbb{P}[X \in A] - \mathbb{P}[X \in A, Y \in A] - \mathbb{P}[X \notin A, Y \in A] \\ &= \mathbb{P}[X \in A, Y \notin A] - \mathbb{P}[X \notin A, Y \in A] \\ &\leq \mathbb{P}[X \in A, Y \notin A]\end{aligned}$$

Coupling and total variation

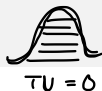
- $TV(\mu, \nu) = \inf\{\mathbb{P}[X \neq Y] : (X, Y) \text{ is a coupling of } \mu \text{ and } \nu\}$.
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- Let (X, Y) be any coupling of μ, ν . Let $A \subseteq \Omega$. Then,

$$\begin{aligned}\mu(A) - \nu(A) &= \mathbb{P}[X \in A] - \mathbb{P}[Y \in A] \\ &= \mathbb{P}[X \in A] - \mathbb{P}[X \in A, Y \in A] - \mathbb{P}[X \notin A, Y \in A] \\ &= \mathbb{P}[X \in A, Y \notin A] - \mathbb{P}[X \notin A, Y \in A] \\ &\leq \mathbb{P}[X \in A, Y \notin A] \\ &\leq \mathbb{P}[X \neq Y].\end{aligned}$$

- The reverse inequality is a starred problem on HW7.

rough idea: $x \in \Omega$, can only get to agree on x
w.p. $\min\{\mu(x), \nu(x)\}$

Poisson approximation



- In our discussion of Poisson random variables, we frequently used the (informal) approximation

$$\text{Pois}(\lambda) \approx X_1 + \dots + X_n, \quad \begin{array}{l} \lambda > 0 \text{ fixed} \\ n \text{ large} \\ (n \rightarrow \infty) \end{array}$$

where X_1, \dots, X_n are i.i.d. Bernoulli random variables with mean λ/n .

- Now, we have the machinery to make this precise.

Poisson approximation

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where X_1, \dots, X_n are i.i.d. Bernoulli random variables with mean λ/n .

- Now, we have the machinery to make this precise.
- Let X_1, \dots, X_n be independent Bernoulli random variables with means p_1, \dots, p_n .
- In other words, for each X_i , $\mathbb{P}[X_i = 1] = p_i$ and $\mathbb{P}[X_i = 0] = (1 - p_i)$.

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- Now, we have the machinery to make this precise.
- Let X_1, \dots, X_n be independent Bernoulli random variables with means p_1, \dots, p_n .
- In other words, for each X_i , $\mathbb{P}[X_i = 1] = p_i$ and $\mathbb{P}[X_i = 0] = (1 - p_i)$.
- Let $S_n = X_1 + \cdots + X_n$.

we are trying to approx
 S_n by $\text{Pois}(\lambda)$

Poisson approximation $\left[\mathbb{P}[\text{Pois}(\lambda_i) = 0] \quad \mathbb{P}[\text{Ber}(p_i) = 0] \right]$

- Let $\lambda_i = -\log(1 - p_i)$. Equivalently, $e^{-\lambda_i} = (1 - p_i)$.
- Let $\lambda = \lambda_1 + \dots + \lambda_n$.

$$S_n = X_1 + \dots + X_n$$

$$X_i \approx \text{Pois}(\lambda_i)$$

$$S_n \approx \text{Pois}(\lambda_1) + \dots + \text{Pois}(\lambda_n)$$

$$\approx \text{Pois}(\lambda)$$

Poisson approximation

- Let $\lambda_i = -\log(1 - p_i)$. Equivalently, $e^{-\lambda_i} = (1 - p_i)$.
- Let $\lambda = \lambda_1 + \cdots + \lambda_n$.
- We will show that

$$\text{TV}(S_n, \text{Pois}(\lambda)) \leq \frac{1}{2} \sum_{i=1}^n \lambda_i^2.$$

Poisson approximation

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$$\text{TV}(S_n, \text{Pois}(\lambda)) \leq \frac{1}{2} \sum_{i=1}^n \lambda_i^2.$$

- **Example:** $\Lambda > 0$ is fixed, n is large, $p_i = \Lambda/n$ for all $i = 1, \dots, n$.
- Then, $\lambda_i = \Lambda/n + O(\Lambda^2/n^2)$, $\lambda = \Lambda + O(\Lambda^2/n)$.

$$-\log\left(1 - \frac{\Lambda}{n}\right)$$

Poisson approximation

- Let $\lambda_i = -\log(1 - p_i)$. Equivalently, $e^{-\lambda_i} = (1 - p_i)$.
- Let $\lambda = \lambda_1 + \dots + \lambda_n$.
- We will show that

$$\text{TV}(S_n, \text{Pois}(\lambda)) \leq \underbrace{\frac{1}{2} \sum_{i=1}^n \lambda_i^2}_{\text{wavy line}}$$

- **Example:** $\Lambda > 0$ is fixed, n is large, $p_i = \Lambda/n$ for all $i = 1, \dots, n$.
- Then, $\lambda_i = \Lambda/n + O(\Lambda^2/n^2)$, $\lambda = \Lambda + O(\Lambda^2/n)$.
- On the homework, you will show that $\text{TV}(\text{Pois}(\mu), \text{Pois}(\nu)) \leq |\nu - \mu|$.

$$\underbrace{\sum_{i=1}^n \lambda_i^2}_{\text{wavy line}} = O\left(\cancel{n} \cdot \frac{\Lambda^2}{n^2}\right) = \underline{\underline{O\left(\frac{\Lambda^2}{n}\right)}}$$

Poisson approximation

- Let $\lambda_i = -\log(1 - p_i)$. Equivalently, $e^{-\lambda_i} = (1 - p_i)$.
- Let $\lambda = \lambda_1 + \dots + \lambda_n$.
- We will show that

$$\text{TV}(S_n, \text{Pois}(\lambda)) \leq \sqrt{\frac{1}{2} \sum_{i=1}^n \lambda_i^2}.$$

- **Example:** $\Lambda > 0$ is fixed, n is large, $p_i = \Lambda/n$ for all $i = 1, \dots, n$.
- Then, $\lambda_i = \Lambda/n + O(\Lambda^2/n^2)$, $\underline{\lambda} = \underline{\Lambda} + \underline{O(\Lambda^2/n)}$.
- On the homework, you will show that $\text{TV}(\text{Pois}(\mu), \text{Pois}(\nu)) \leq |\nu - \mu|$.
- Then, by the triangle inequality,

$$\text{TV}(S_n, \text{Pois}(\Lambda)) \leq \text{TV}(S_n, \text{Pois}(\lambda)) + \text{TV}(\text{Pois}(\lambda), \text{Pois}(\Lambda))$$

S_n is connected to $\text{Pois}(\lambda)$ and $\text{Pois}(\Lambda)$ by arrows.

 The first term is $O\left(\frac{\Lambda^2}{n}\right)$.

 The second term is $O\left(\frac{\Lambda^2}{n}\right)$.

Poisson approximation

- Let $\bar{\lambda}_i = -\log(1 - p_i)$. Equivalently, $e^{-\lambda_i} = (1 - p_i)$.
- Let $\lambda = \lambda_1 + \dots + \lambda_n$.
- We will show that

$$\text{TV}(S_n, \text{Pois}(\lambda)) \leq \frac{1}{2} \sum_{i=1}^n \lambda_i^2.$$

- **Example:** $\Lambda > 0$ is fixed, n is large, $p_i = \Lambda/n$ for all $i = 1, \dots, n$.
- Then, $\bar{\lambda}_i = \Lambda/n + O(\Lambda^2/n^2)$, $\lambda = \Lambda + O(\Lambda^2/n)$. $\lambda_i = -\log(1 - p_i)$
- On the homework, you will show that $\text{TV}(\text{Pois}(\mu), \text{Pois}(\nu)) \leq |\nu - \mu|$.
- Then, by the triangle inequality,

$$\begin{aligned} \text{TV}(S_n, \text{Pois}(\Lambda)) &\leq \text{TV}(S_n, \text{Pois}(\lambda)) + \text{TV}(\text{Pois}(\lambda), \text{Pois}(\Lambda)) \\ &\leq O(\Lambda^2/n) + O(\Lambda^2/n) \\ &\leq O(\Lambda^2/n), \end{aligned}$$

which justifies our approximation from before.

Poisson approximation

- We now show that

$$\text{TV}(S_n, \text{Pois}(\lambda)) \leq \frac{1}{2} \sum_{i=1}^n \lambda_i^2.$$

let me try to prove this for $n=1$.

$$\left[\text{TV}(\text{Ber}(p), \text{Pois}(-\log(1-p))) \right] \leq \frac{1}{2} (\log(1-p))^2$$

Poisson approximation

- We now show that

$$\text{TV}(S_n, \text{Pois}(\lambda)) \leq \frac{1}{2} \sum_{i=1}^n \lambda_i^2.$$

- Let us first prove this for $n = 1$. Let $\lambda = -\log(1 - p)$. We want to show:

$$\text{TV}(\text{Ber}(p), \text{Pois}(\lambda)) \leq \frac{1}{2} \lambda^2.$$

Poisson approximation

- We now show that

$$\text{TV}(S_n, \text{Pois}(\lambda)) \leq \frac{1}{2} \sum_{i=1}^n \lambda_i^2.$$

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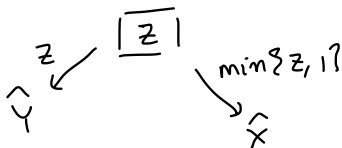
$$\text{TV}(\text{Ber}(p), \text{Pois}(\lambda)) \leq \frac{1}{2} \lambda^2.$$

- By the coupling lemma, it suffices to exhibit a coupling (\hat{X}, \hat{Y}) of $\text{Ber}(p)$ and $\text{Pois}(\lambda)$ such that

$$\mathbb{P}[\hat{X} \neq \hat{Y}] \leq \frac{1}{2} \lambda^2.$$

Poisson approximation

- Here is such a coupling: Generate $Z \sim \text{Pois}(\lambda)$. Then, set $\hat{Y} = Z$ and $\hat{X} = \min\{Z, 1\}$.



Poisson approximation

- Here is such a coupling: Generate $Z \sim \text{Pois}(\lambda)$. Then, set $\hat{Y} = Z$ and $\hat{X} = \min\{Z, 1\}$.
- Clearly \hat{Y} has the correct marginal distribution. As for \hat{X} , note that

$$\begin{aligned} \mathbb{P}[\hat{X} = 0] &= \mathbb{P}[Z = 0] = e^{-\lambda} = (1 - p) = \mathbb{P}[\text{Ber}(p) = 0]. \\ \Rightarrow \mathbb{P}[\hat{X} = 1] &= 1 - (1 - p) = p \end{aligned}$$

$$\hat{X} = \min\{Z, 1\} \quad \hat{X} = 0 \Leftrightarrow Z = 0$$

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$$\mathbb{P}[\hat{X} = 0] = \mathbb{P}[Z = 0] = e^{-\lambda} = (1 - p) = \mathbb{P}[\text{Ber}(p) = 0].$$

- Moreover,

$$\mathbb{P}[\hat{X} \neq \hat{Y}] = \mathbb{P}[Z \geq 2]$$

$$\begin{aligned}\hat{Y} &= Z \\ \hat{X} &= \min\{Z, 1\}\end{aligned}$$

Poisson approximation

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$$\mathbb{P}[\hat{X} \neq \hat{Y}] = \mathbb{P}[Z \geq 2]$$

$$= e^{-\lambda} \sum_{j \geq 2} \frac{\lambda^j}{j!}$$

$$= e^{-\lambda} \frac{\lambda^2}{2} \sum_{j \geq 2} \frac{\lambda^{j-2}}{j!} \cdot 2$$

change of
vars.
 $j-2 = k$

Poisson approximation

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- Clearly \hat{Y} has the correct marginal distribution. As for \hat{X} , note that

$$\mathbb{P}[\hat{X} = 0] = \mathbb{P}[Z = 0] = e^{-\lambda} = (1 - p) = \mathbb{P}[\text{Ber}(p) = 0].$$

- Moreover,

$$\begin{aligned}\mathbb{P}[\hat{X} \neq \hat{Y}] &= \mathbb{P}[Z \geq 2] \\ &= e^{-\lambda} \sum_{j \geq 2} \frac{\lambda^j}{j!} \\ &\leq \frac{\lambda^2}{2} \underbrace{\sum_{j \geq 0} e^{-\lambda} \frac{\lambda^j}{j!}}_1\end{aligned}$$

Poisson approximation

- Here is such a coupling: Generate $Z \sim \text{Pois}(\lambda)$. Then, set $\hat{Y} = Z$ and $\hat{X} = \min\{Z, 1\}$.
- Clearly \hat{Y} has the correct marginal distribution. As for \hat{X} , note that

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- Moreover,

$$\begin{aligned}\mathbb{P}[\hat{X} \neq \hat{Y}] &= \mathbb{P}[Z \geq 2] \\ &= e^{-\lambda} \sum_{j \geq 2} \frac{\lambda^j}{j!} \\ &\leq \frac{\lambda^2}{2} \sum_{j \geq 0} e^{-\lambda} \frac{\lambda^j}{j!} \quad \lambda = -\log(1-p). \\ &= \frac{\lambda^2}{2}. \quad \Rightarrow \text{TV}(\text{Ber}(p), \text{Pois}(\lambda)) \leq \lambda^2/2\end{aligned}$$

Poisson approximation

- At this point, we are almost done.

Poisson approximation

$$S_n = X_1 + \dots + X_n \text{ ind.}$$
$$\text{Poi}(\lambda) = \text{Poi}(\lambda_1) + \dots + \text{Poi}(\lambda_n)$$

independent

- At this point, we are almost done.
- Let (\hat{X}_i, \hat{Y}_i) be a coupling of $\text{Ber}(p_i)$ and $\text{Pois}(\lambda_i)$ as above.
- Let $(\hat{X}_1, \hat{Y}_1), \dots, (\hat{X}_n, \hat{Y}_n)$ be independent copies of this coupling.

$$Z_1 \sim \text{Poi}(\lambda_1) \qquad Z_n \sim \text{Poi}(\lambda_n) \qquad Z_1 \dots Z_n \text{ are ind.}$$
$$\hat{Y}_1 = Z_1$$
$$\hat{X}_1 = \min\{Z_1, 1\} \quad \dots \quad \dots$$

Poisson approximation

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- Then, $S_n \sim \widehat{X}_1 + \dots + \widehat{X}_n$ and $\text{Pois}(\lambda) \sim \widehat{Y}_1 + \dots + \widehat{Y}_n$.

$$\begin{array}{ccc} \text{"} & & \text{"} \\ X_1 + \dots + X_n & & \text{Pois}(\lambda_1) + \dots + \text{Pois}(\lambda_n) \end{array}$$

•

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- Let $(\hat{X}_1, \hat{Y}_1), \dots, (\hat{X}_n, \hat{Y}_n)$ be independent copies of this coupling.
- Then, $S_n \sim \hat{X}_1 + \dots + \hat{X}_n$ and $\text{Pois}(\lambda) \sim \hat{Y}_1 + \dots + \hat{Y}_n$.
- Moreover, by the coupling lemma and our previous calculation

$$\text{TV}(S_n, \text{Pois}(\lambda)) \leq \mathbb{P}[\hat{X}_1 + \dots + \hat{X}_n \neq \hat{Y}_1 + \dots + \hat{Y}_n]$$

Poisson approximation

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- Moreover, by the coupling lemma and our previous calculation

$$\begin{aligned}\text{TV}(S_n, \text{Pois}(\lambda)) &\leq \mathbb{P}[\widehat{X}_1 + \dots + \widehat{X}_n \neq \widehat{Y}_1 + \dots + \widehat{Y}_n] \\ &\leq \mathbb{P}[\widehat{X}_1 \neq \widehat{Y}_1] + \dots + \mathbb{P}[\widehat{X}_n \neq \widehat{Y}_n] \\ &\leq \frac{1}{2} \lambda_1^2 + \dots + \frac{1}{2} \lambda_n^2\end{aligned}$$

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- Moreover, by the coupling lemma and our previous calculation

$$\begin{aligned}\text{TV}(S_n, \text{Pois}(\lambda)) &\leq \mathbb{P}[\hat{X}_1 + \dots + \hat{X}_n \neq \hat{Y}_1 + \dots + \hat{Y}_n] \\ &\leq \mathbb{P}[\hat{X}_1 \neq \hat{Y}_1] + \dots + \mathbb{P}[\hat{X}_n \neq \hat{Y}_n] \\ &\leq \frac{\lambda_1^2}{2} + \dots + \frac{\lambda_n^2}{2}.\end{aligned}$$

$\forall A$

$$\mathbb{P}_r(S_n \in A) = \mathbb{P}_r(\text{Pois}(\lambda) \in A) \pm \left(\frac{\lambda_1^2 + \dots + \lambda_n^2}{2} \right)$$

Random mapping representation of Markov chains

- We have often specified transitions of Markov chains in words. For instance, for the symmetric simple random walk, instead of writing down the transition matrix, we have used a simple description like: at each step, toss an independent fair coin. If the coin lands heads, move one step right. Else, move one step left.

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- We can formalize this by using the **random mapping representation** of a transition matrix P on the state space S . This is simply a function $f : S \times \Lambda \rightarrow S$ along with a Λ -valued random variable Z which satisfies

$$\mathbb{P}[f(x, Z) = y] = P_{x,y}.$$

- For instance, in the case of the symmetric simple random walk, we can take $\Lambda = \{\underline{H}, \underline{T}\}$, Z is a random variable which is H with probability $1/2$ and T otherwise, and $\underline{f(x, H)} = \underline{x + 1}$, $\underline{f(x, T)} = \underline{x - 1}$.

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- Indeed, we can take $\Lambda = [0, 1]$, Z is uniformly distributed on $[0, 1]$ and

$$f(i, z) = j \iff \sum_{\ell=1}^{j-1} P_{i,\ell} \leq z \leq \sum_{\ell=1}^j P_{i,\ell}.$$

