

STATS 217: Introduction to Stochastic Processes I

Lecture 2

Recall from last time

- Given integers $A > 0, B > 0$, let

$$\tau := \min\{n \geq 0 : S_n = A \text{ or } S_n = -B\}.$$

- For $-B \leq k \leq A$, define

$$g(k) := \mathbb{E}[\tau \mid S_0 = k].$$

- Clearly, $g(-B) = 0, g(A) = 0$.

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- Clearly, $\tilde{g}(-B) = 0, \tilde{g}(A) = 0$.

- For $-B < k < A$, we have

$$\begin{aligned} g(k) &= \frac{1}{2} \mathbb{E}[\tau \mid S_0 = k, X_1 = 1] + \frac{1}{2} \mathbb{E}[\tau \mid S_0 = k, X_1 = -1] \\ &= \frac{1}{2} (g(k+1) + 1) + \frac{1}{2} (g(k-1) + 1) \\ &= \frac{1}{2} \underset{\sim}{g}(k+1) + \frac{1}{2} \underset{\sim}{g}(k-1) + \underset{\sim}{1}. \end{aligned}$$

First step analysis

- Let $(\Delta h)(k) := h(k+1) - h(k)$.
- Then, for all $-B < k < A$

$$\begin{aligned}(\Delta(\Delta g))(k-1) &= (\Delta g)(k) - (\Delta g)(k-1) \\ &= g(k+1) - g(k) - g(k) + g(k-1) \\ &= g(k+1) - (g(k+1) + g(k-1) + 2) + g(k-1) \\ &= -2.\end{aligned}$$

- “Second derivative of g is -2 ” so $g(k) = -k^2 + Dk + C$.
- Using boundary conditions,

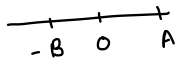
$$g(k) = -(k-A)(k+B).$$

First step analysis

Therefore,

$$g(k) = \mathbb{E}[\tau \mid S_0 = k] = -(k \cancel{+} A)(k + B).$$

in our example



$$\tau_{(A, -B)}$$

$$g(0) = AB$$

First step analysis

Therefore,

$$g(k) = \mathbb{E}[\tau \mid S_0 = k] = -\frac{\bar{c}}{(k \wedge A)(k + B)}.$$

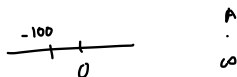
- **Answer 3:** $A = 200, B = 100, g(0) = 2 \times 10^4$.

First step analysis

Therefore,

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- **Answer 4 (ii):** " $A = \infty$ ", $B = 100, g(0) = \infty$.



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- Formally, let

$$\tau_1 = \min\{n \geq 0 : S_n = -100\},$$

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$$\begin{aligned}\tau_1 &= \min\{n \geq 0 : S_n = -100\}, \\ \tau_2(\ell) &= \min\{n \geq 0 : S_n = -100 \text{ or } S_n = \ell\} \quad \forall \ell \geq 1.\end{aligned}$$

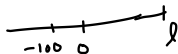
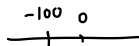
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$$\tau_2(\ell) = \min\{n \geq 0 : S_n = -100 \text{ or } S_n = \ell\} \quad \forall \ell \geq 1.$$

- Then, for all $\ell \geq 1$, $\tau_2(\ell) \leq \tau_1$ so that

$$g(0) = \frac{-(-100)(0+1)}{100\ell} = \mathbb{E}[\tau_2(\ell) \mid S_0 = 0] \leq \mathbb{E}[\tau_1 \mid S_0 = 0],$$

and now take $\ell \rightarrow \infty$.

$\tau_2(\ell) \mid S_0 = 0 \leq \tau_1 \mid S_0 = 0$

First step analysis

-

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- On the other hand, **Answer 4(i)**: $\tau_2(\ell) = \text{min time to hit } -100 \text{ or } +\ell$

$$\begin{aligned}\mathbb{P}[S_n \text{ visits } -100] &\geq \mathbb{P}[S_{\tau_2(\ell)} = -100] \\ &= \frac{\ell}{100 + \ell}\end{aligned}$$

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- So, a symmetric simple random walk starting at 0 visits -100 with probability 1. Again, there is nothing special about -100 here.

Summary

We have studied some aspects of the Gambler's Ruin.

- What is the probability that a symmetric simple random walk started from 0 hits 2 before -1 ? We saw that this is $1/3$.
- What is the expectation of the first time when the walk hits either 2 or -1 ? We saw that this is 2.
- Moreover, we saw that while the probability of hitting 1 is 1, the expectation of the first time we hit 1 is infinite.

$$\mathbb{E}[\tau_1] = \sum_{k=1}^{\infty} \underbrace{\mathbb{P}[\tau_1 = k]}_{\text{wavy}} \cdot k$$

Path counting and applications

Today, we will develop tools that allow us to answer questions like the following:

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- What is the probability that the maximum value of the first 2020 steps of the random walk is 10?

Path counting and applications

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- What is the probability that the first time we hit 1 is exactly 101 steps?
- What is the probability that the random walk stays non-negative for the first 2020 steps?
- What is the probability that the maximum value of the first 2020 steps of the random walk is 10?
- ...and more!

Path counting

We will need the following notation:

- $N_n(a, b)$ = number of paths from a to b with n steps.
(wavy underline under n)

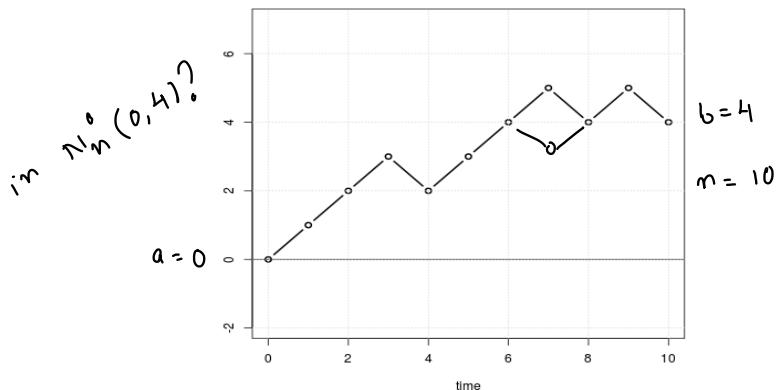


Image courtesy www.isical.ac.in

Path counting

at time $1, \dots, n-1$

- $N_n^0(a, b)$ = number of paths from a to b with n steps that visit 0 ~~after time λ~~ ⁰.
- $N_n^{\neq 0}(a, b)$ = number of paths from a to b with n steps that do not visit 0 at times $1, 2, \dots, n-1$.

Path counting

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- $N_n^{\neq 0}(a, b)$ = number of paths from a to b with n steps that do not visit 0 at times $1, 2, \dots, n-1$.

Note the following direct consequences of the definitions.

- $N_n(a, b) = N_n^{\neq 0}(a, b) + N_n^0(a, b)$.

Path counting

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Note the following direct consequences of the definitions.

- $N_n(a, b) = N_n^{\neq 0}(a, b) + N_n^0(a, b)$.

- Also, $N_n(a, b) = N_n^0(a, b)$ if a and b have different signs.

*easy
to compute*

$$a = -2 \quad b = 5$$

Path counting

Let us compute $N_n(a, b)$.

- Let u denote the number of $+1$ steps and d denote the number of -1 steps.
- Since the path has n steps, we must have $u + d = n$.

Path counting

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- Let u denote the number of $+1$ steps and d denote the number of -1 steps.
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- Since the path goes from a to b , we must have $u - d = b - a$.

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- Since the path has n steps, we must have $u + d = n$.
- Since the path goes from a to b , we must have $u - d = b - a$.
- Hence, $u = (n + b - a)/2$ so that

$$\frac{+1}{-} \quad - \quad - \quad - \quad \frac{+1}{-} \quad - \quad \frac{+1}{-}$$

$$u = 3 \\ n = 7$$

n steps
 u of them are $+1$

Path counting

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- Since the path goes from a to b , we must have $u - d = b - a$.
- Hence, $u = (n + b - a)/2$ so that

$$N_n(a, b) = \binom{n}{(n+b-a)/2}, \quad \binom{n}{k} = \frac{n!}{k!(n-k)!}$$

with the convention that $\binom{n}{r} = 0$ if r is not an integer.

Reflection principle

For any $a > 0$ and $b > 0$,

- $N_n^0(a, b) = N_n(-a, b)$.

and this
we can
compute

how can we compute
 $N_n^0(a, b)$? say $a \neq 0, b \neq 0$
① a and b have opp. signs.
you know how to do this

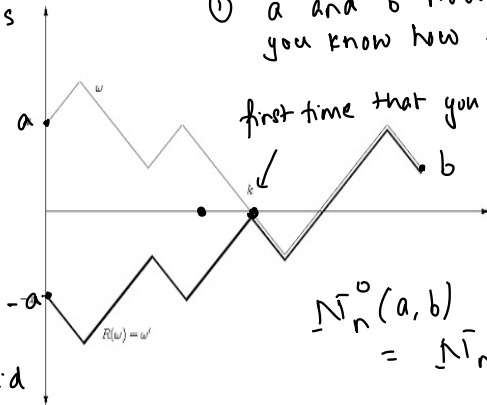
$$N_n^0(a, b)$$

$$f \downarrow \uparrow g$$

$$N_n(-a, b)$$

$$f \circ g = \text{id}$$

$$g \circ f = \text{id}$$



$$N_n^0(a, b) = N_n(-a, b)$$

Image courtesy www.tricki.org

Reflection principle

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- So, $N_n^{\neq 0}(a, b) = N_n(a, b) - N_n(-a, b)$.

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- So, $N_n^{\neq 0}(a, b) = N_n(a, b) - N_n(-a, b)$.

The point is that we already have a formula for the expressions on the right hand side.

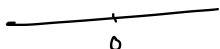
Return time to 0

- Let $(S_n)_{n \geq 0}$ be a simple, symmetric random walk starting from 0.
- Let $\tau_0 := \inf\{n \geq 1 : S_n = 0\}$.
- What is the pmf of τ_0 ?

$$\mathbb{P}[\tau_0 = k] \quad \forall k \in \mathbb{N}^+$$

$$\mathbb{P}[\tau_0 = 1] = ? \quad \mathbb{P}[\tau_0 = 3] = ?$$

$$\mathbb{P}[\tau_0 = \text{odd number}] = 0$$



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- Observe that the support of τ_0 consists of even natural numbers.

values k s.t. $\mathbb{P}[\tau_0 = k] \neq 0$

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- Let $\tau_0 := \inf\{n \geq 1 : S_n = 0\}$.
- What is the pmf of τ_0 ?
- Observe that the support of τ_0 consists of even natural numbers.
- Moreover, for any $k \geq 1$

$$\mathbb{P}[\tau_0 = 2k] = \underbrace{N_{2k}^{\neq 0}(0,0)}_{\substack{\text{# of fav. paths} \\ \text{~~~~~}}} \cdot 2^{-2k}.$$

- start at 0
- end at 0
- take $2k$ steps
- do not hit 0 in the middle.

2^{2k} =
number of
paths of
length $2k$.

Return time to 0

To compute $N_{2k}^{\neq 0}(0,0)$, we can use the reflection principle.

$$N_{2k}^{\neq 0}(0,0) \stackrel{\text{first step analysis}}{=} \underbrace{N_{2k-1}^{\neq 0}(1,0)} + \underbrace{N_{2k-1}^{\neq 0}(-1,0)}$$
$$N_{2k-1}^{\neq 0}(1,0) > N_{2k-1}^{\neq 0}(1,0)$$

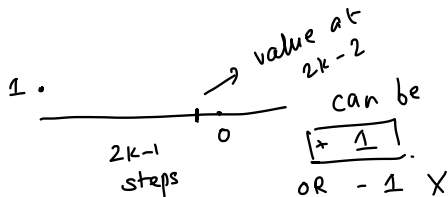
?

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$$\begin{aligned} N_{2k}^{\neq 0}(0,0) &= N_{2k-1}^{\neq 0}(1,0) + N_{2k-1}^{\neq 0}(-1,0) \\ &= 2N_{2k-1}^{\neq 0}(1,0) \end{aligned}$$

$$\begin{aligned} N_{2k-1}^{\neq 0}(1,0) &= N_{2k-2}^{\neq 0}(1,1) \end{aligned}$$



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Return time to 0

- Simplifying the arithmetic, we get that

$$N_{2k}^{\neq 0}(0, 0) = \frac{1}{2k-1} \binom{2k}{k}.$$

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- Hence,

$$\begin{aligned} \mathbb{P}[\tau_0 = 2k] &= \frac{1}{2k-1} \binom{2k}{k} 2^{-2k} \\ &= \frac{1}{2k-1} \mathbb{P}[S_{2k} = 0]. \end{aligned}$$

The Ballot Problem

- Consider an election with two candidates A and B .
- Suppose that a votes have been cast for A and b votes have been cast for B where $a > b$.
- After the votes have been cast, they are counted in a uniformly random order.

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- Suppose that a votes have been cast for A and b votes have been cast for B where $a > b$.
- After the votes have been cast, they are counted in a uniformly random order.
- Since $a > b$, after all the votes are counted, A emerges as the winner.
- What is the probability that A leads B throughout the count?

$$\underbrace{\# \text{ votes for } A}_{\text{counted}} > \underbrace{\# \text{ votes for } B}_{\text{counted}}$$

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$$\mathbb{P}[A \text{ leads throughout}] = \frac{N_{a+b}^{\neq 0}(0, a - b)}{N_{a+b}(0, a - b)}.$$

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- So, it only remains to compute $N_{a+b}^{\neq 0}(0, a - b)$.

The Ballot Problem

We need to compute $N_{a+b}^{\neq 0}(0, a - b)$.

$$\begin{aligned} N_{a+b}^{\neq 0}(0, a - b) &= N_{a+b-1}^{\neq 0}(1, a - b) \\ &= N_{a+b-1}(1, a - b) - N_{a+b-1}^0(1, a - b) \end{aligned}$$

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Hence,

$$\mathbb{P}[A \text{ leads throughout}] = \frac{a - b}{a + b}.$$

The Ballot Problem

- One way to reinterpret the conclusion of the Ballot problem is that for any $a > b \geq 0$ and for a simple symmetric random walk starting from $S_0 = 0$,

$$\mathbb{P}[S_i > 0 \quad \forall i = 1, \dots, a + b - 1 \mid S_{a+b} = a - b] = \frac{a - b}{a + b}.$$

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- Rewritten in more convenient notation, for any integers $k, n > 0$,

$$\mathbb{P}[S_1 > 0, \dots, S_{n-1} > 0, S_n = k] = \frac{k}{n} \cdot \mathbb{P}[S_n = k].$$

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- On the homework, you will explore variants of this.